Course introduction

Advanced Compiler Construction
Michel Schinz – 2013-02-21

General information

Course goal

The goal of this course is to teach you:

- how to compile high-level functional and object-oriented programming languages,
- how to optimize the generated code, and
- how to support code execution at run time.

To achieve these goals, the course is roughly split in three parts of unequal length:

- a part covering the compilation of high-level concepts (e.g. closures),
- a part covering intermediate languages and optimizations,
- a part covering virtual machines and garbage collection.

Evaluation

The grade will be based on three aspects:

- two group projects, to be completed in groups of at most two people,
- two individual projects, to be completed alone,
- an individual oral exam.

Warning: the course is evaluated during the semester, which has two important consequences:

- there is no retake exam,
- the oral exam will take place during the last week of the semester, not during the official exam period.

Grading scheme

The final grade will be based on your results in:

- the various project parts, spread over 11 weeks, which contribute to 80% of the grade,
- the final exam, which contributes to 20% of the grade.

Resources

Lecturer:
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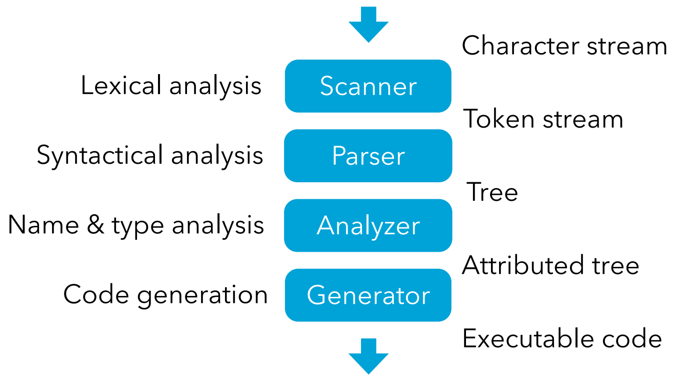
Assistant:
  Vlad Ureche, INR 320, `Vlad.Ureche@epfl.ch`
Web page:
  `http://lamp.epfl.ch/teaching/advanced_compiler`
Moodle site:
  `http://moodle.epfl.ch/course/view.php?id=13857`

Course overview

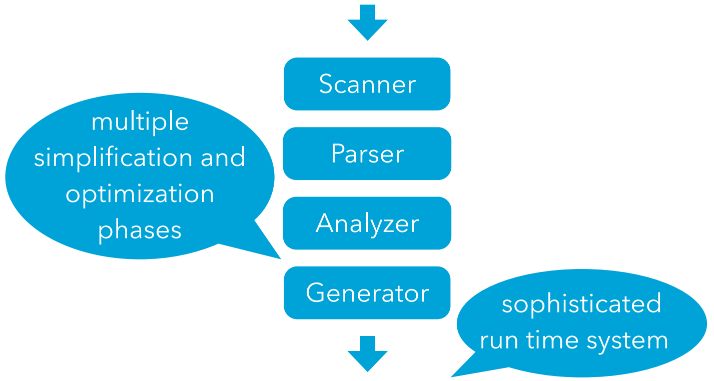
What is a compiler

Your current view of a compiler must be something like this:



What is a compiler, really?

Real compilers are often more complicated…


Scanner
Parser
Analyzer
Generator
multiple simplification and optimization phases
sophisticated run time system

Additional phases

Simplification phases transform the program so that complex concepts of the language – pattern matching, anonymous functions, etc. – are translated using simpler ones.

Optimization phases transform the program so that it hopefully makes better use of some resource – e.g. CPU cycles, memory, etc.

Of course, all these phases must preserve the meaning of the original program!

Simplification phases

Example of a simplification phase: Java compilers have a phase that transforms nested classes to top-level ones.

```
class Out {
  void f1() { }
  class In {
    void f2() {
      f1();
    }
  }
}
```

→

```
class Out {
  void f1() { }
}
class Out$In {
  final Out this$0;
  Out$In(Out o) {
    this$0 = o;
  }
  void f2() {
    this$0.f1();
  }
}
```

Optimization phases

Example of an optimization phase: Java compilers optimize expressions involving constant values. That includes removing **dead code**, i.e. code that can never be executed.

```
class C {
  public final static boolean debug = !true;
  int f() {
    if (debug) {
      System.out.println("C.f() called");
    }
    return 10;
  }
}
```

dead code, removed during compilation

Intermediate representations

To manipulate the program, simplification and optimization phases must represent it in some way.
One possibility is to use the representation produced by the parser – the abstract syntax tree (AST).
The AST is perfectly suited to certain tasks, but other **intermediate representations** (**IR**) exist and are more appropriate in some situations.

Run time system

Implementing a high-level programming language usually means more than just writing a compiler. A complete **run time system** must be written, to provide various services to executing programs, like:

- code loading and linking,
- code interpretations, compilation and optimization,
- memory management (garbage collection),
- concurrency,
- etc.

This is quite a lot, and modern Java Virtual Machines, for example, are more complex than Java compilers!

Memory management

Most modern programming languages offer **automatic memory management**: the programmer allocates memory explicitly, but deallocation is performed automatically.
The deallocation of memory is usually performed by a part of the run time system called the **garbage collector** (**GC**).
A garbage collector periodically frees all memory that has been allocated by the program but is not reachable anymore.

Virtual machines

Instead of targeting a real processor, a compiler can target a virtual one, usually called a **virtual machine** (**VM**). The produced code is then interpreted by a program emulating the virtual machine.
Virtual machines have many advantages:

- the compiler can target a single architecture,
- the program can easily be monitored during execution, e.g. to prevent malicious behavior, or provide debugging facilities,
- the distribution of compiled code is easier.

The main (only?) disadvantage of virtual machines is their speed: it is always slower to interpret a program in software than to execute it directly in hardware.

Dynamic (JIT) compilation

To make virtual machines faster, **dynamic**, or **just-in-time** (**JIT**) compilation was invented.
The idea is simple: Instead of interpreting a piece of code, the virtual machine translates it to machine code, and hands that code to the processor for execution.
This is usually faster than interpretation.

Summary

Compilers for high-level languages are more complex than the ones you've studied, since:

- they must translate high-level concepts like pattern-matching, anonymous functions, etc. to lower-level equivalents,
- they must be accompanied by a sophisticated run time system, and
- they should produce optimized code.

This course will be focused on these aspects of compilers and run time systems.