Currying

Motivation

Look again at the summation functions:

```
def sumInts(a: Int, b: Int) = sum(x => x, a, b)
def sumCubes(a: Int, b: Int) = sum(x => x * x * x, a, b)
def sumFactorials(a: Int, b: Int) = sum(fact, a, b)
```

Question

Note that a and b get passed unchanged from sumInts and sumCubes into sum.

Can we be even shorter by getting rid of these parameters?

Functions Returning Functions

Let's rewrite sum as follows.

```
def sum(f: Int => Int): (Int, Int) => Int = {
    def sumF(a: Int, b: Int): Int =
        if (a > b) 0
        else f(a) + sumF(a + 1, b)
        sumF
}
```

sum is now a function that returns another function.

The returned function sumF applies the given function parameter f and sums the results.

Stepwise Applications

We can then define:

def sumInts = sum(x => x)
def sumCubes = sum(x => x * x * x)
def sumFactorials = sum(fact)

These functions can in turn be applied like any other function:

```
sumCubes(1, 10) + sumFactorials(10, 20)
```

Consecutive Stepwise Applications

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- sum(cube) is therefore equivalent to sumCubes.
- ▶ This function is next applied to the arguments (1, 10).

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Generally, function application associates to the left:

sum(cube)(1, 10) == (sum (cube)) (1, 10)

Multiple Parameter Lists

The definition of functions that return functions is so useful in functional programming that there is a special syntax for it in Scala.

For example, the following definition of sum is equivalent to the one with the nested sumF function, but shorter:

```
def sum(f: Int => Int)(a: Int, b: Int): Int =
    if (a > b) 0 else f(a) + sum(f)(a + 1, b)
```

Expansion of Multiple Parameter Lists

In general, a definition of a function with multiple parameter lists

```
\text{def } f(\text{args}_1)...(\text{args}_n) = E
```

where n > 1, is equivalent to

 $def \ f(args_1)...(args_{n-1}) = \{def \ g(args_n) = E; g\}$

where g is a fresh identifier. Or for short:

 $\text{def } f(\text{args}_1)...(\text{args}_{n-1}) = (\text{args}_n \Rightarrow E)$

Expansion of Multiple Parameter Lists (2)

By repeating the process n times

 $\text{def } f(\text{args}_1)...(\text{args}_{n-1})(\text{args}_n) = E$

is shown to be equivalent to

def $f = (args_1 \Rightarrow (args_2 \Rightarrow ...(args_n \Rightarrow E)...))$

This style of definition and function application is called *currying*, named for its instigator, Haskell Brooks Curry (1900-1982), a twentieth century logician.

In fact, the idea goes back even further to Schönfinkel and Frege, but the term "currying" has stuck.

More Function Types

Question: Given,

```
def sum(f: Int => Int)(a: Int, b: Int): Int = ...
```

What is the type of sum ?

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Answer:

```
(Int => Int) => (Int, Int) => Int
```

Note that functional types associate to the right. That is to say that

Int => Int => Int

is equivalent to

Int => (Int => Int)

Exercise

The sum function uses linear recursion. Write a tail-recursive version by replacing the ???s.

```
def sum(f: Int => Int)(a: Int, b: Int): Int = {
    def loop(a: Int, acc: Int): Int = {
        if (???) ???
        else loop(???, ???)
    }
    loop(???, ???)
}
```